## Parallel Performance of Structured AMR Calculations Using the SAMRAI Framework

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with

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August 14, 2002 BlueGene/L Workshop





### **Outline**

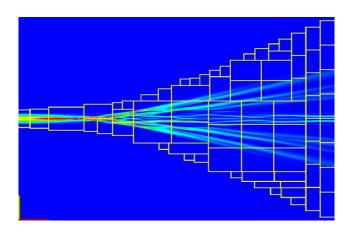
- SAMRAI introduction
- Parallel implementation of SAMR
- Parallel performance measurements
- New algorithms to enhance parallel performance
- Requirements and issues on BG/L

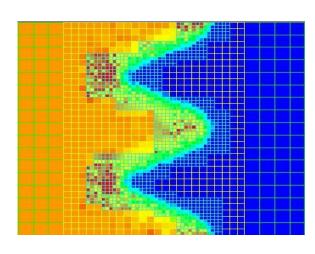
### SAMRAI

Structured Adaptive Mesh Refinement Application Infrastructure

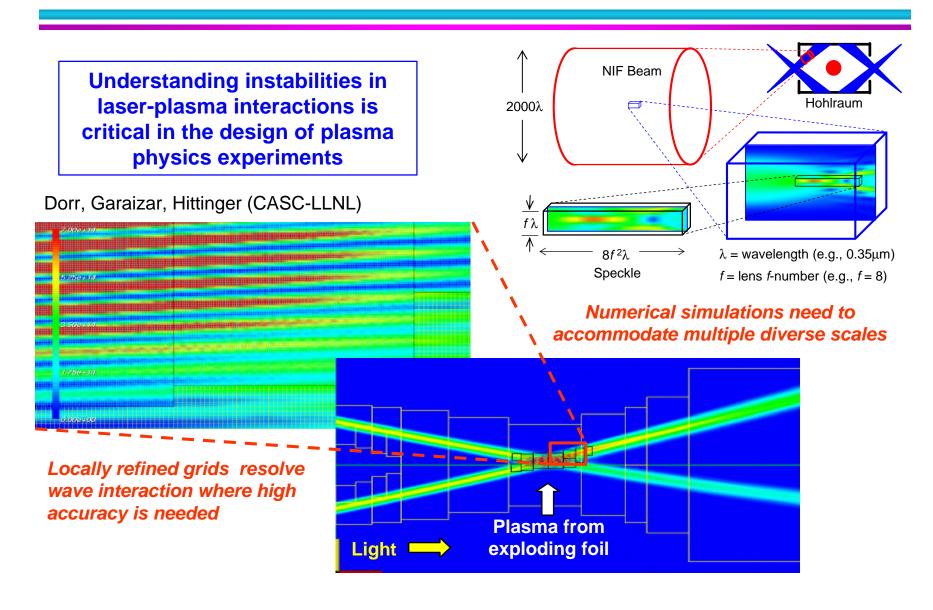
- Object-oriented (C++) software framework for parallel (MPI) adaptive multi-physics applications
- Supports applications investigating multi-scale phenomena.
- High-level reusable code and algorithms shared across a variety of applications.

www.llnl.gov/CASC/SAMRAI





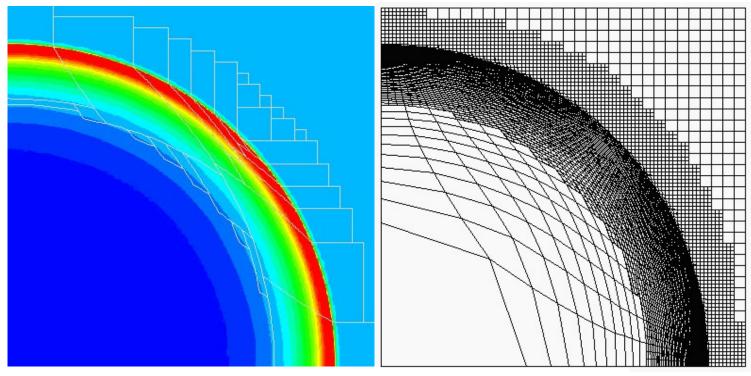
# ALPS uses SAMRAI for adaptive laser plasma instability simulation



# ALE-AMR couples ALE models with AMR to model shock hydrodynamics

Improve accuracy of ALE simulations by increasing concentration of mesh points around regions of interest

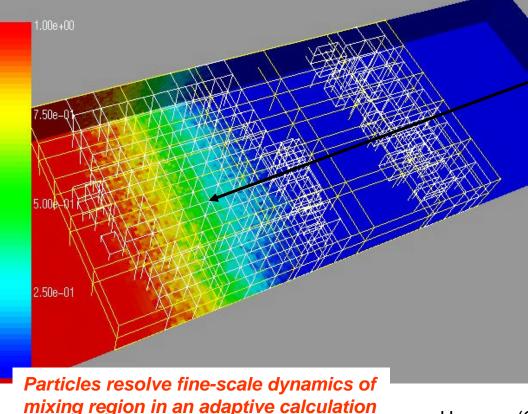
Anderson, Pember, Elliott (CASC-LLNL)



Sedov blast wave density and Lagrangian mesh

# Hybrid continuum-DSMC model used to efficiently resolve interface dynamics

Interface instability problems (e.g., Richtmyer-Meshkov) involve coarse-scale hydrodynamic transport and fine-scale molecular diffusion



fluid A fluid B

Continuum representation
(Euler, Navier-Stokes)
away from interface

DSMC representation at interface

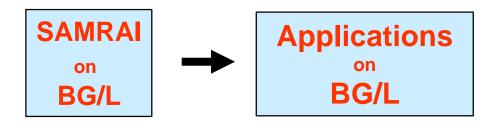
- Interface region grows and moves as instability evolves
- Standard CFD simulation of turbulent mixing is limited by finest mesh scale
- Particle resolve molecular behavior but are too expensive for large domains

Hornung (CASC), Garcia (SJSU)

shock

## SAMRAI provides infrastructure support for a variety of applications

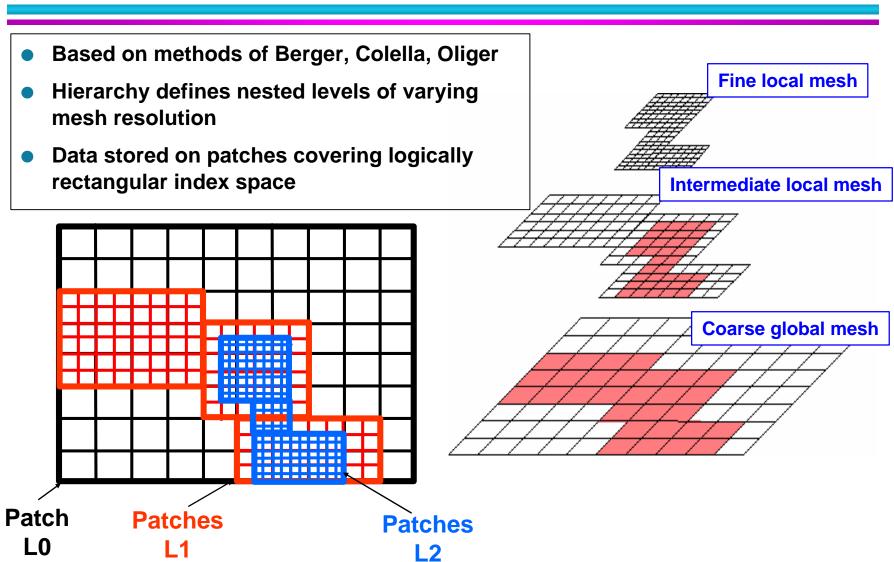
- Parallel processing support (MPI)
- Shared algorithms
- Interfaces for SAMR data to solvers (PETSc, PVODE, hypre)
- Checkpointing & restart support (HDF)
- Parallel tools (VAMPIR, TAU)
- Current users regularly run on existing large processor systems



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# Structured AMR (SAMR) employs a dynamically adaptive "patch" hierarchy



### Dynamic mesh adapts to features as solution evolves

### Adaptive solution of **Euler equations**

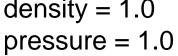
### **Initial conditions:**

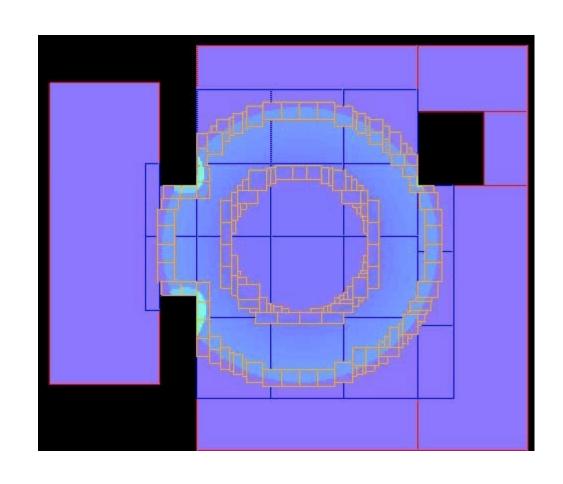
inside sphere

density = 8.0pressure = 40.0

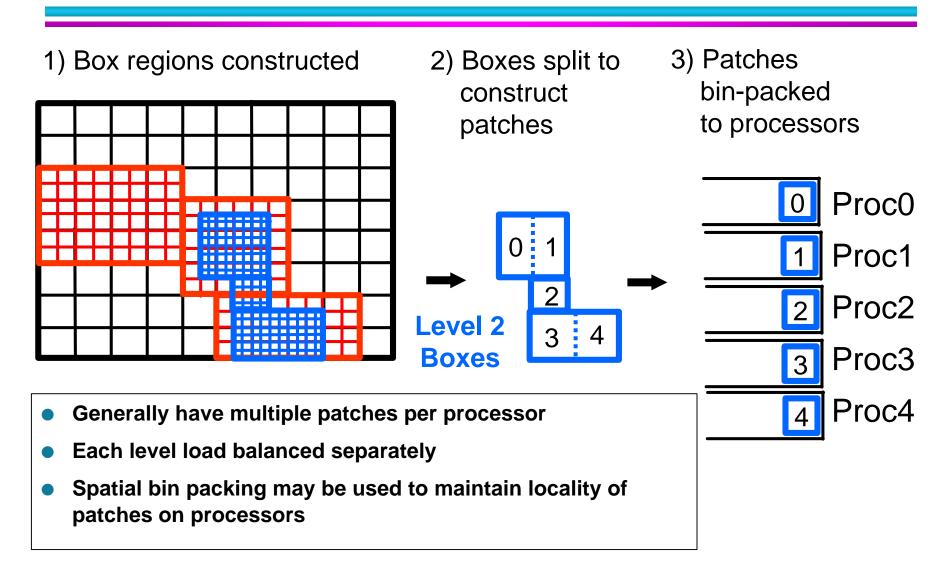
### outside sphere

density = 1.0



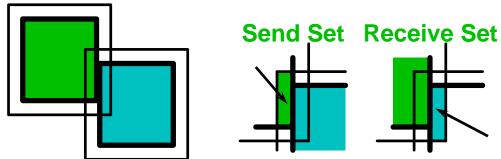


## Patches distributed to processors to balance computational workload

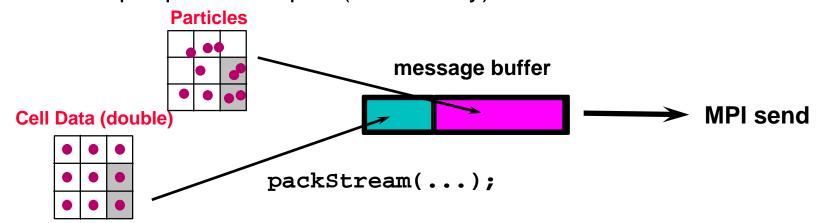


## Communication schedules create and store data dependencies

 Amortize cost of creating send/receive sets over multiple communication cycles



- Data from various sources packed into single message stream
  - supports complicated variable-length data
  - one send per processor pair (low latency)



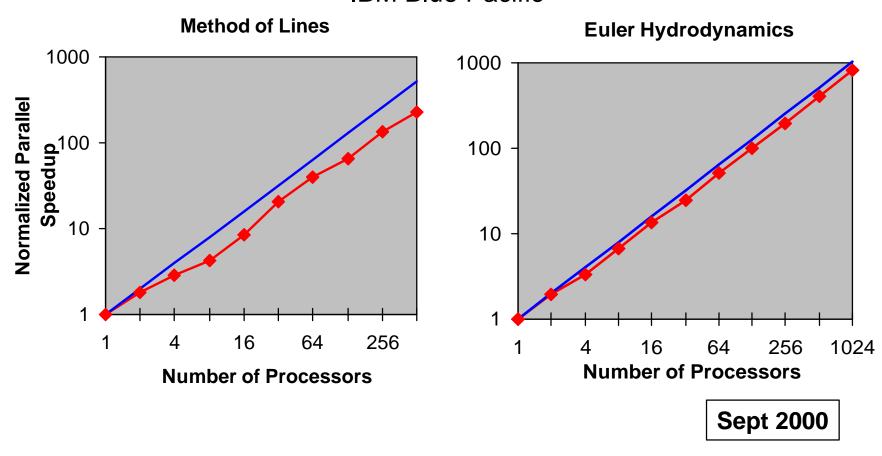
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# Non-adaptive calculations using SAMRAI show good scaling

### **Single Level calculation**

50x50x50 patch per processor IBM Blue Pacific



## Benchmarks constructed to analyze scaling properties of SAMR applications

- Simple numerical kernels
- Invoke the main algorithmic components used in more complex apps (e.g. meshing, time advance, etc.)
- Timing decomposed into two phases:

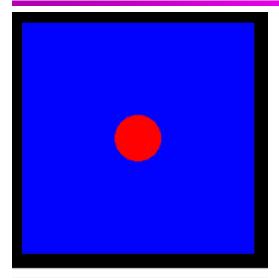
#### **Time Advance:**

- numerical kernels
- communication (filling ghost cells)
- load imbalances

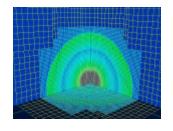
### **Re-Gridding:**

- cluster tagged cells
- construct communication schedules
- distributing data to new mesh configuration

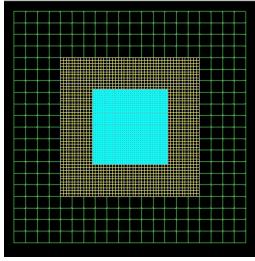
# Non-scaled Euler benchmark – same problem size run on all processors

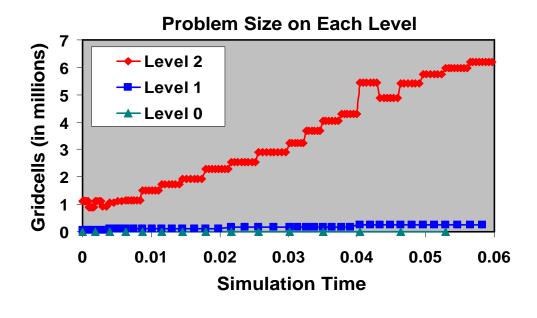


### 3D spherical shock - Euler hydrodynamics

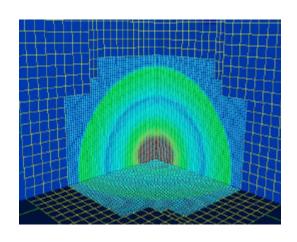


- Workload changes over simulation
- Per-processor workload decreases as number of processors increased



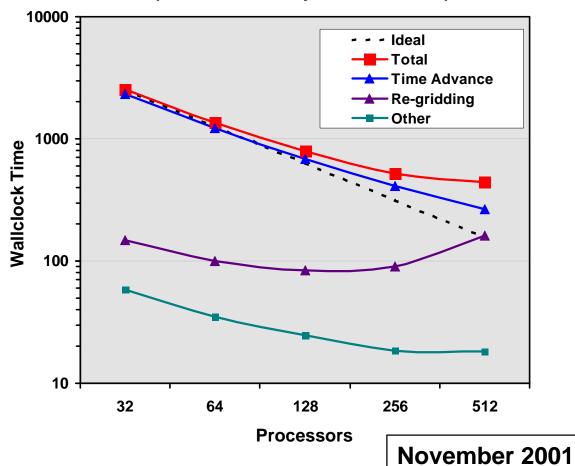


## Parallel Performance of *non-scaled* adaptive Euler benchmark

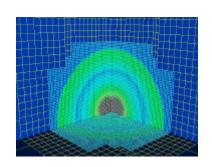


Non-scaled
Euler calculation
ASCI IBM Blue Pacific

### Measured Solution Time on Various Processors (3 Level Euler Sphere Problem)



# Poor scaling in re-gridding hurts efficiency on large processor counts

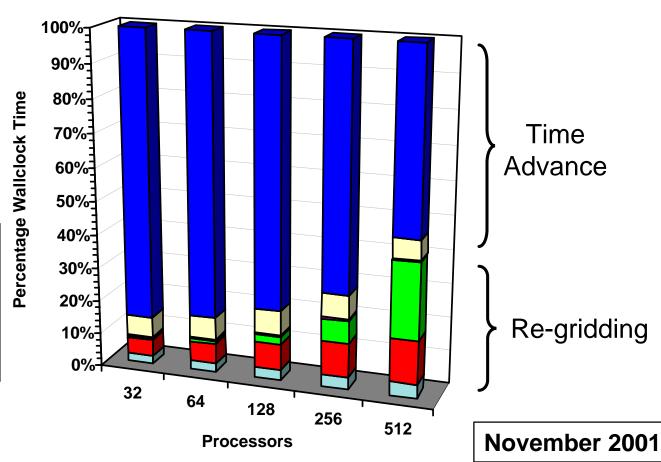


Non-scaled Euler calculation

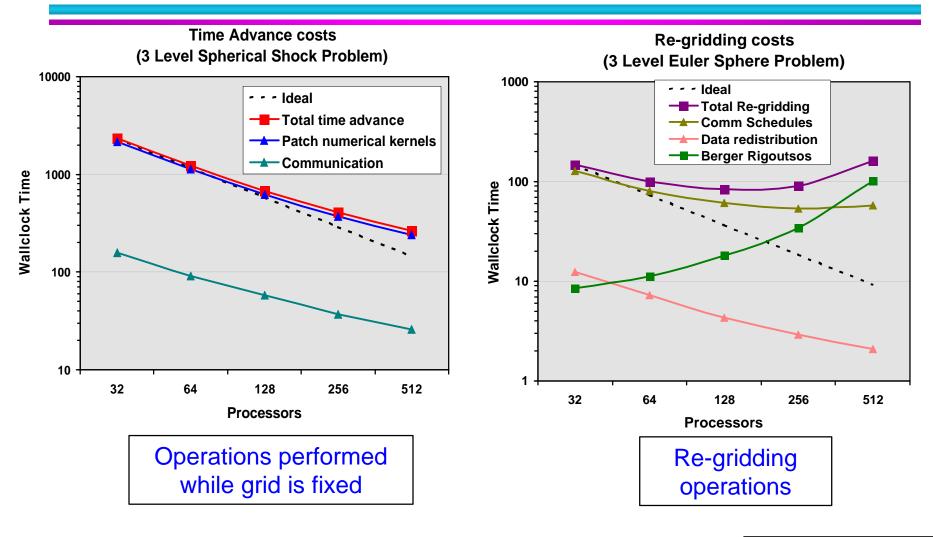
ASCI IBM Blue Pacific



Measured Solution Time on Various Processors (3 Level Spherical Shock Problem)

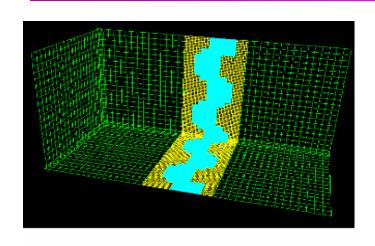


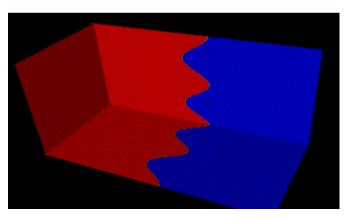
# Poor scaling in re-gridding hurts efficiency on large processor counts (ASCI Blue Pac)



**November 2001** 

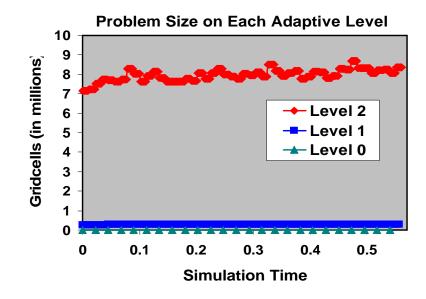
## Scaled linear advection benchmark – problem size increased with processors



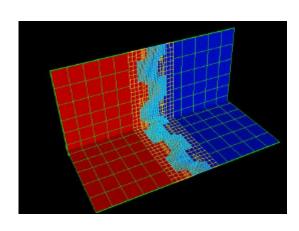


#### 3D advecting sinusoidal front - linear advection

- Workload uniform over simulation
- Per-processor workload remains constant as number of processors is increased

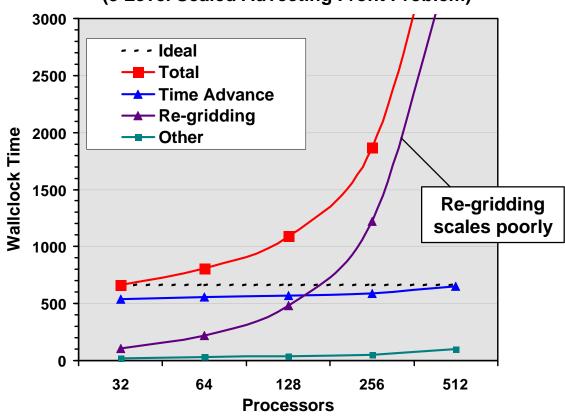


### Parallel performance of scaled linear advection benchmark



Scaled
Linear advection
calculation
ASCI IBM Blue Pacific

### Measured Solution Time on Various Processors (3 Level Scaled Advecting Front Problem)



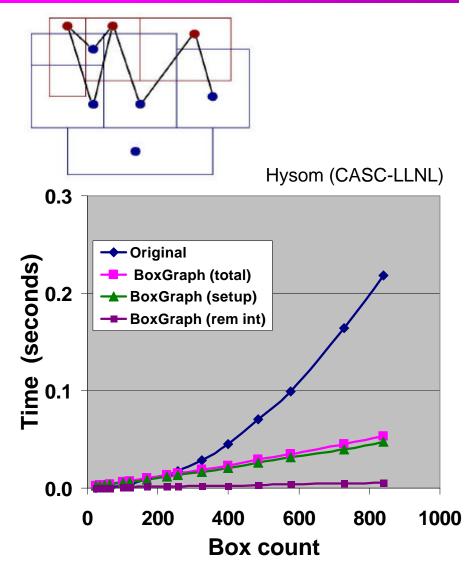
November 2001

### **Outline**

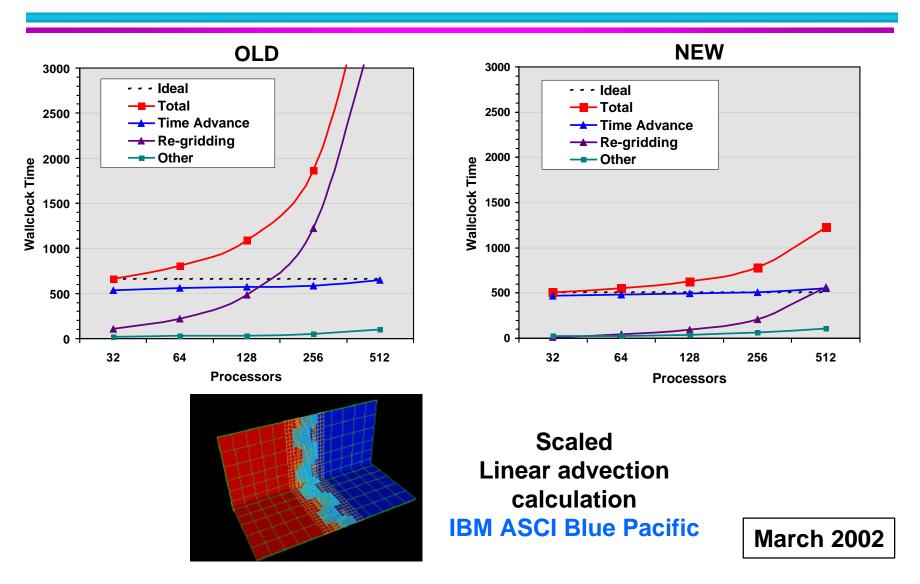
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### Graph-based algorithms speed up communication schedule construction

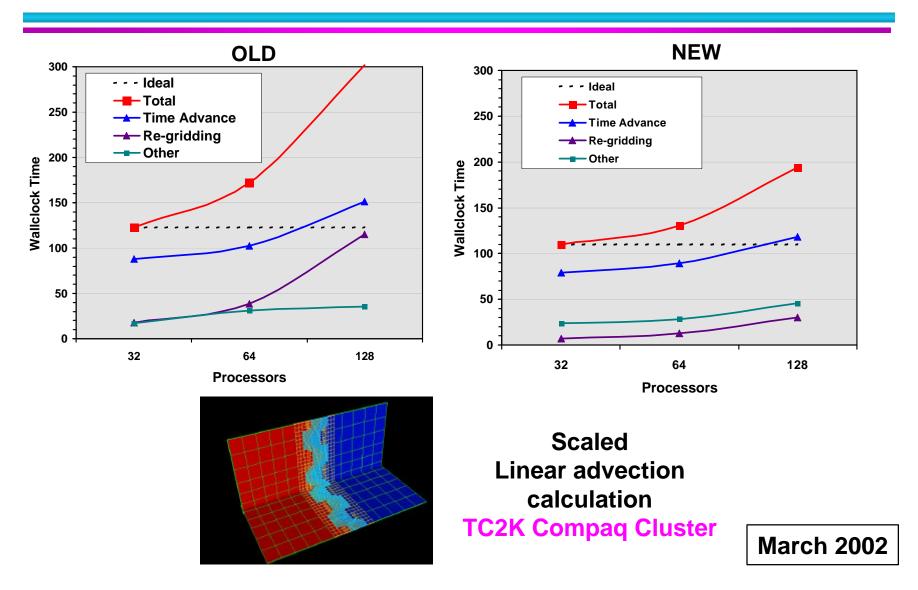
- Before constructing a communication schedule to transfer data between two levels, build a "Box graph":
  - Insert a vertex in V for each box
  - Insert an edge (i,j) at intersection
- Given this graph, can find a box's neighbors in O(1)
- Primary cost is graph construction



# Scaled results with new graph-based schedule construction algorithm



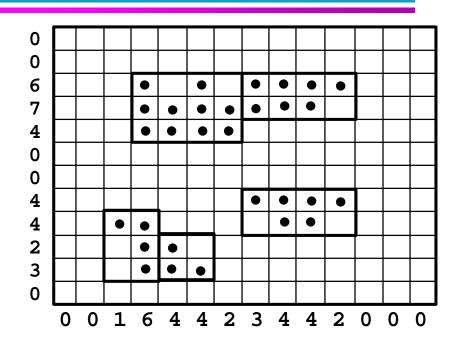
# Scaled results with new graph-based schedule construction algorithm



# Binary Tree reduction for tagged-cell clustering algorithm (Berger-Rigoutsos)

### Berger-Rigoutsos:

- Forms new patches from tagged cells
- Determines box cuts from global histogram through recursive algorithm
- Original implementation used global reductions to form histogram

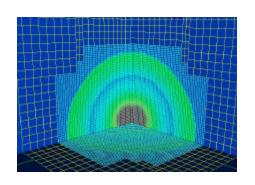


### New Implementation:

- Binary-tree reduction algorithm collects information from selected processors at each recursion.
- New box configuration constructed and broadcast by one processor.

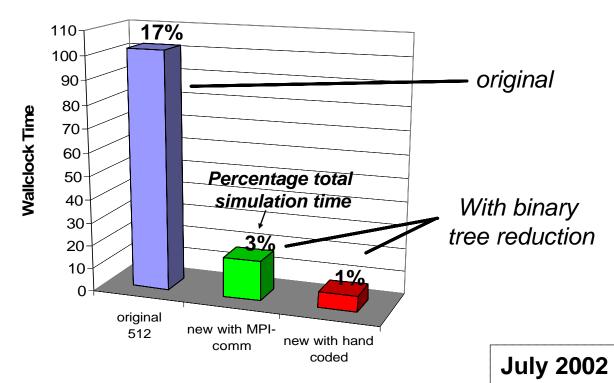
## Timing results Berger-Rigoutsos algorithm with binary-tree reductions.

- Binary tree reduction algorithm two implementations
  - MPI communicators
  - Hand-coded MPI send/recvs



Non-Scaled
Euler calculation
IBM ASCI Blue Pacific

Berger-Rigoutsos – 512 processors



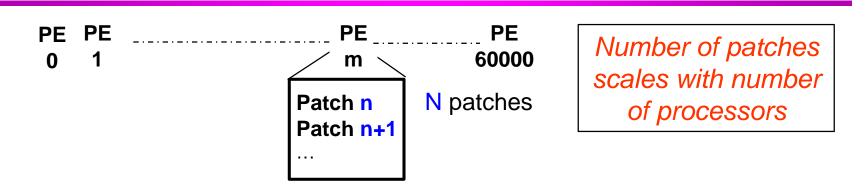
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### BlueGene/L wish list

- SAMRAI Dependencies:
  - C++, C, F77/F90 compilers
  - MPI
  - HDF5 (checkpointing)
- Desirable features:
  - C++-capable debugging tool
  - Memory analysis tool (i.e. reports stack/heap usage on nodes)

# Scaling issues with a large number of processors



- Box operations in gridding may invoke O(N²) algorithms (e.g. former communication schedule algorithm).
- More efficient graph-based algorithms work on up to 512 processors, but need to develop efficient algorithms for O(10K) processors.
- Difficult to assess beforehand because complexity is generally problem dependent.

# BG/L will require us to rethink our grid storage approach

 Current approach: Each MPI process holds a box for each patch in the problem to determine communication dependencies.

 Because # patches grows with # processors, trivial overhead becomes non-trivial on BG/L.

		per processor	
procs	<u>patches</u>	storage (MB)	Large overhead for nodes of BG/L
0.5K 60K	2.5K-10K 300K-1200K	< 1 MB 20-80MB	

### Collective communications on BG/L

- Berger-Rigoutsos clustering:
  - Binary tree reduction algorithm effective in reducing costs on O(0.5K) processors.
  - Will this approach be effective on O(10K) processors?
- Some global communications are necessary (e.g. timestep synchronization in time advance).

### **Concluding Remarks**

- Porting SAMRAI to BG/L enables a variety of applications to use the architecture.
- Results of scaling AMR algorithms on up to 512 processors:
  - Communication not the primary source of scaling inefficiency.
  - Re-gridding operations that are trivial on small numbers of processors become significant on large numbers.
  - More efficient graph-based algorithms successful in reducing these costs.
- Speculation on running AMR applications on BG/L:
  - Re-gridding costs will likely be the main hindrance.
  - Continued exploration into more efficient gridding algorithms needed.

### **Auspices Statement**

- This work was performed under the auspices of the U.S.
   Department of Energy by University of California Lawrence
   Livermore National Laboratory under contract No. W-7405-Eng-48.
- Document UCRL-PRES-149437